

Logics for Sizes with Union or Intersection

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What Kind of Inference?

$\frac{\text{All cats are mammals that purr} \quad \text{There are at least as many cats as purring things}}{\text{All purring things are cats}}$

$\frac{\text{All hippies can vote} \quad \text{All hippies are vegetarian} \quad \text{All people can vote} \quad \text{There are at least as many people as things that can vote}}{\text{There are at least as many hippies as people}} \quad \frac{\text{All vegetarian people are hippies}}{\text{All vegetarian people are hippies}}$

Semantics

$[a] \subseteq M$
 $[a \cap b] = [a] \cap [b]$
 $\mathcal{M} \models \text{AtLeast } x \ y \text{ iff } |[x]| \geq |[y]|$
 $\mathcal{M} \models \text{All } x \ y \text{ iff } [x] \subseteq [y]$
 $\Gamma \models \varphi$

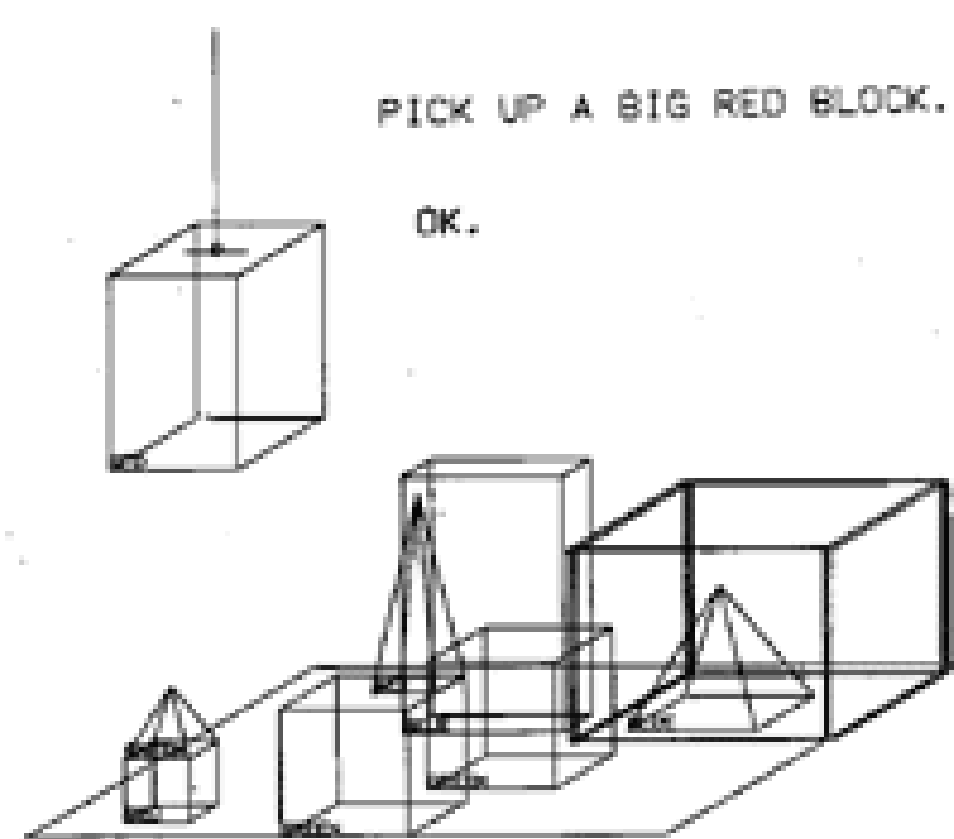
Main Technical Results

Theorem. *The logics $\mathcal{A}^\cap(\text{card})$ and $\mathcal{A}^\cup(\text{card})$ are complete.*

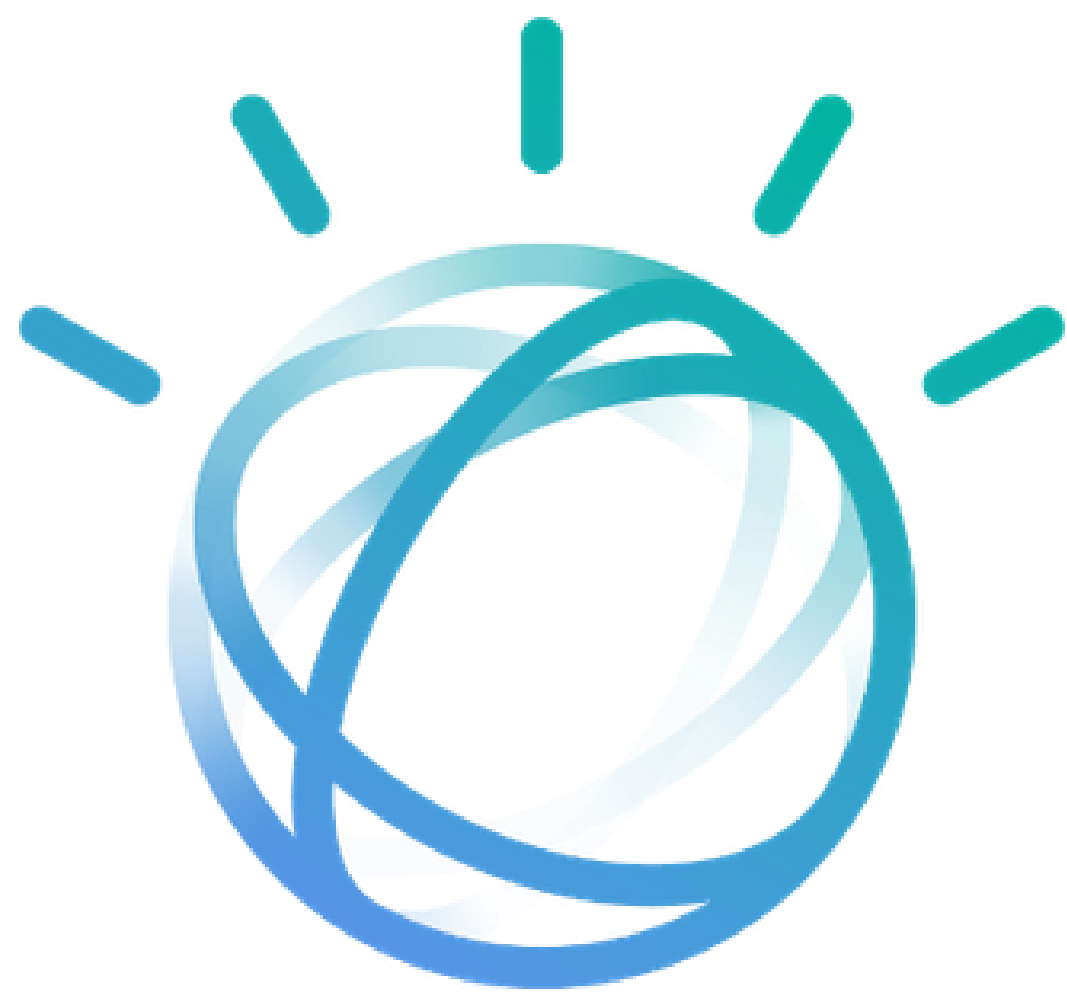
Theorem. *\vdash is decidable in polynomial time!*

Theorem. *If $\Gamma \not\models \varphi$, then we can construct a *finite* countermodel \mathcal{M} satisfying Γ but falsifying φ in polynomial time*

Examples in AI



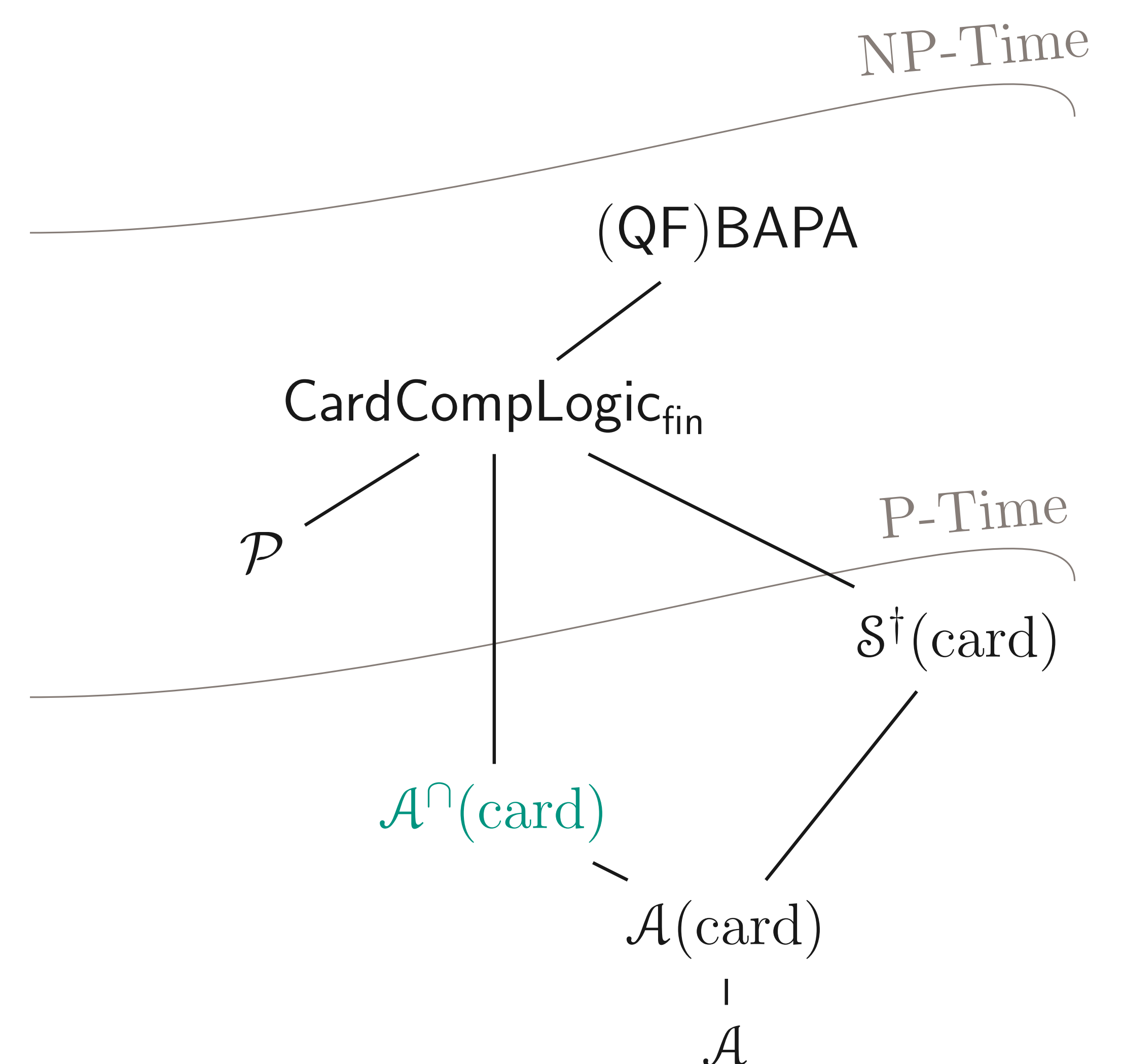
Z3



Inference Rules

$\frac{}{\text{All } x \ x} \text{ (AXIOM)}$	$\frac{\text{All } x \ y \quad \text{All } y \ z}{\text{All } x \ z} \text{ (BARBARA)}$	
$\frac{\text{All } x \ y \quad \text{AtLeast } x \ y}{\text{All } y \ x} \text{ (MIX)}$	$\frac{\text{All } x \ y}{\text{AtLeast } y \ x} \text{ (SIZE)}$	$\frac{\text{AtLeast } x \ y \quad \text{AtLeast } y \ z}{\text{AtLeast } x \ z} \text{ (TRANS)}$
$\frac{}{\text{All } (a \cap b) \ a} \text{ (INTER-L)}$	$\frac{}{\text{All } (a \cap b) \ b} \text{ (INTER-R)}$	$\frac{\text{All } a \ b \quad \text{All } a \ c}{\text{All } a \ (b \cap c)} \text{ (INTER-ALL)}$
$\frac{}{\text{All } a \ (a \cup b)} \text{ (UNION-L)}$	$\frac{}{\text{All } b \ (a \cup b)} \text{ (UNION-R)}$	$\frac{\text{All } a \ c \quad \text{All } b \ c}{\text{All } (a \cup b) \ c} \text{ (UNION-ALL)}$
$\frac{\text{More } x \ y \quad \text{AtLeast } y \ z}{\text{More } x \ z} \text{ (MORE-L)}$	$\frac{\text{AtLeast } x \ y \quad \text{More } y \ z}{\text{More } x \ z} \text{ (MORE-R)}$	
$\frac{\text{More } x \ y}{\text{AtLeast } x \ y} \text{ (MORE-ATLEAST)}$	$\frac{\text{More } z \ z}{\varphi} \text{ (X)}$	$\frac{\text{[AtLeast } y \ x]}{\text{More } z \ z} \text{ (RAA)}$

The Map



Future Work

Polynomial Time with 'More'

$$\frac{\text{All } x \ a \quad \text{All } x \ b \quad \text{More } (a \cup b) \ b}{\text{More } a \ x}$$

Adding 'Some'

$$\frac{\text{More } a \ b \quad \text{AtLeast } c \ d \quad \text{AtLeast } (b \cup d) \ (a \cup c)}{\text{Some } a \ c}$$

Integrating with Z3 Plugin

- Our logic may best be viewed as an efficient fragment of (QF)BAPA
- Consequently, we keep an eye towards embedding $\mathcal{A}^\cap(\text{card})$ in an SMT solver
- Ideal candidate is Z3 plugin [6] for BAPA

References

- [1] Yifeng Ding et al. "The Logic of Comparative Cardinality". 2018.
- [2] Viktor Kuncak and Martin Rinard. "Towards efficient satisfiability checking for boolean algebra with Presburger Arithmetic". 2007.
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- [5] Ian Pratt-Hartmann. "On the Computational Complexity of the Numerically Definite Syllogistic and Related Logics". 2008.
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